Lecture 1 The Role of Algorithms in Computing

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Chapter 1 Topics

- Algorithms
- Algorithms as a Technology

Goals:

Learn techniques of algorithm design and analysis so that you can:

- develop algorithms,
- show that they give the correct answer, and
- understand their efficiency

What is an algorithm?

• An *algorithm* is any well-defined computational procedure that takes some value, or set of values, as input and produces some value, or set of values, as output. An algorithm is thus a sequence of computational steps that transform the *input* into the *output*. (CLRS, p. 5)

What is an algorithm?

• An *algorithm* is a tool for solving a well-specified computational problem. The statement of the problem specifies in general terms the desired input/output relationship . The algorithm describes a specific computational procedure for achieving that input/output relationship. (CLRS, p. 5)

Formal definition of the *sorting problem*:

Input: A sequence of numbers $\langle a_1, a_2, K, a_n \rangle$

Output: A permutation (reordering) $\langle a'_1, a'_2, K, a'_n \rangle$ of the input sequence such that $a'_1 \le a'_2 \le K \le a'_n$

Instance: The input sequence <14, 2, 9, 6, 3> is an *instance* of the sorting problem.

An *instance* of a problem consists of the input (satisfying whatever constraints are imposed in the problem statement) needed to compute a solution to the problem.

Correctness: An algorithm is said to be *correct* if, for every instance, it halts with the correct output. We say that a correct algorithm *solves* the given computational problem.

Efficiency: Algorithms that solve the same problem can differ enormously in their efficiency. Generally speaking, we would like to select the most efficient algorithm for solving a given problem.

Time efficiency: When we talk about the efficiency of an algorithm, we usually mean the *time* requirements of the algorithm: how long would it take a program executing this algorithm to solve the problem?

Space efficiency: Space efficiency is usually an all-ornothing proposition; either we have enough space in our computer's memory to run a program implementing a specific algorithm, or we do not. If we have enough, we're OK; if not, we can't run the program at all. Consequently, analysis of the space requirements of a program tend to be pretty simple and straightforward.

Space efficiency:

- Note that space requirements set a minimum lower bound on the time efficiency of the problem.
- Suppose that our data structure is a single-dimensioned array with n = 100 elements in it. Let's say that the first step in our algorithm is to execute a loop that just copies values into each of the 100 elements. Then our algorithm must take at least 100 iterations of the loop. So the running time of our algorithm is at least O(n), just from setting up (initializing) our data structure!

Time efficiency of two sorts:

Suppose we use insertion sort to sort a list of numbers. Insertion sort has a time efficiency roughly equivalent to $c_1 \cdot n^2$. The value n is the number of items to be sorted. The value c_1 is a constant which represents the overhead involved in running this algorithm; it is independent of n.

Compare this to merge sort. Merge sort has a time efficiency of $c_2 \cdot n \cdot lg n$ (where lg n is the same as log_2n).

n	Insertion sort - $O(n^{2})$	Merge sort – $O(n lg n)$
4	16	8
8	64	24
16	256	64
32	1024	160
64	4096	384
128	16,394	896
256	65,536	2048
512	262,144	4608
1024	1,048,576	10,240
1,048,576	~1,000,000,000,000	20,971,520

Time efficiency of two sorts:

Do the two constants, c_1 and c_2 , affect the result?

Yes, but only with low values of n.

- Suppose that insertion sort is hand-coded in machine language for optimal performance and its overhead is very low, so that $c_1 = 2$.
- Now suppose that merge sort is written in Ada by an average programmer and the compiler doesn't do a good job of optimization, so that $c_2 = 50$.

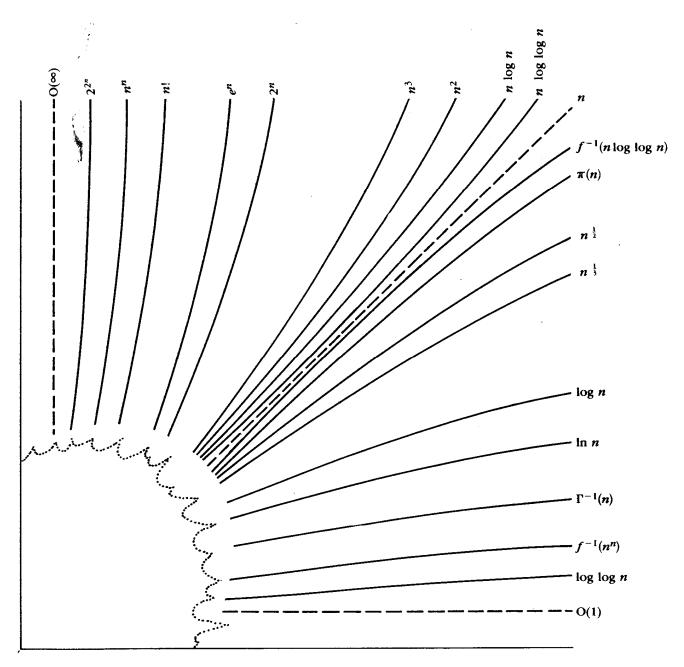
Time efficiency of two sorts:

To make things worse, suppose that insertion sort is run on a machine that executes 1 billion instructions per second, while merge sort is run on a slow machine that executes only 10 million instructions per second.

Now let's sort 1 million numbers:

insertion sort: $\frac{2 \cdot (10^6)^2 instructions}{10^9 instructions / sec} = 2000 sec.$

merge sort: $\frac{50 \cdot 10^6 \, \text{lg} 10^6 \, \text{instructions}}{10^7 \, \text{instructions / sec}} \approx 100 \, \text{sec.}$



From E. D. Reilly and F. D. Federighi, *Pascalgorithms*, Boston: Houghton Mifflin, 1989, p. 501.

Figure 14.3. Relative asymptotic behavior of typical f(n) (not drawn to scale)

Conclusion

What does this mean for computer science?

It means that using efficient algorithms can be even more important that building faster computers: more efficient thinking beats more efficient hardware!

And this means that algorithms are definitely worth studying.